

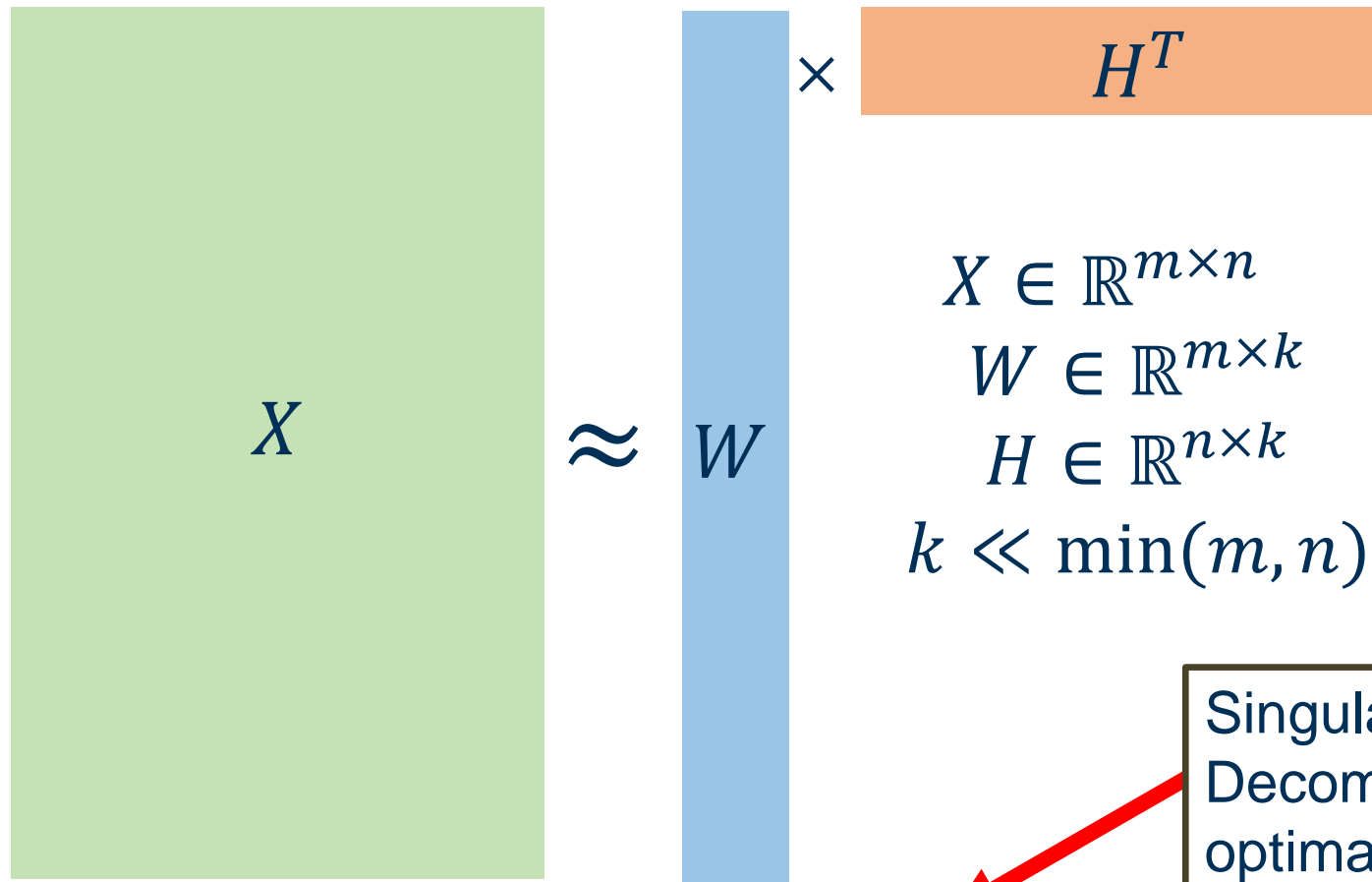
# Constrained Low-rank Approximation

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1. What is Low-rank Approximation?
2. Applications of Low-rank Approximation
3. Randomized Algorithms for Low-rank Approximation

# Low-rank Matrix Approximation



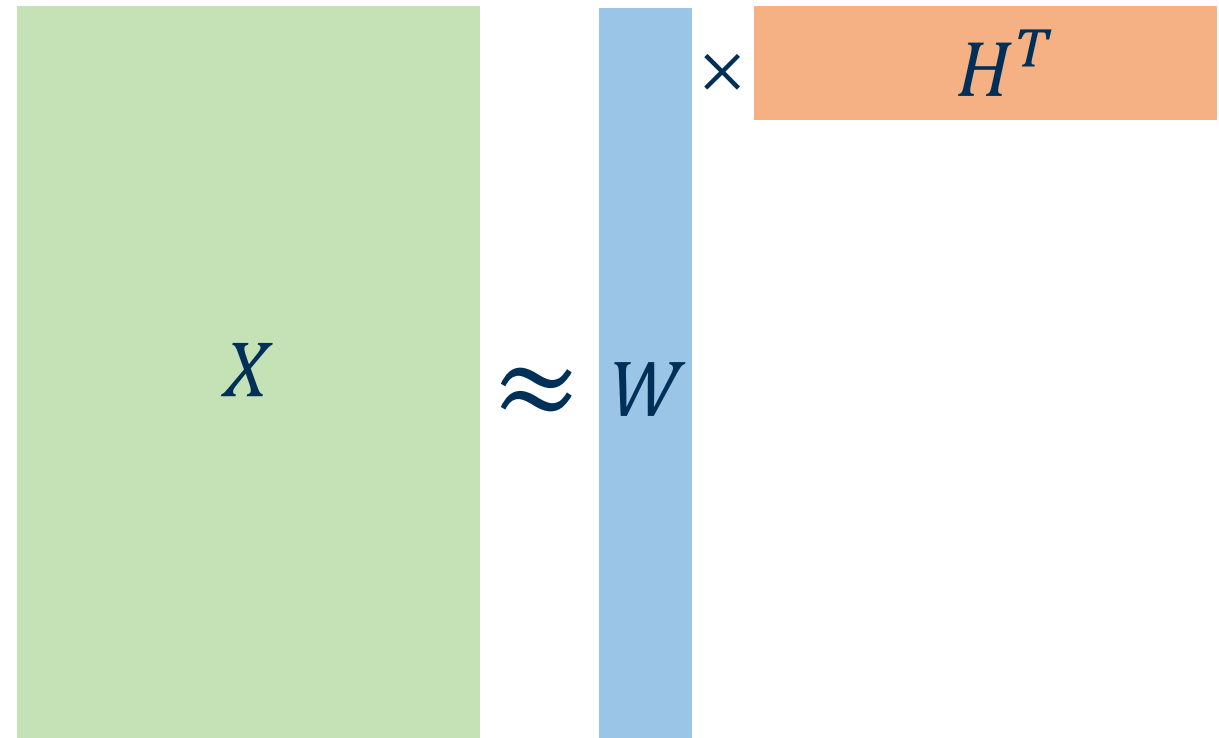
$$\min_{\{W, H\}} \|X - WH^T\|_F$$

# Nonnegative Matrix Factorization (NMF)

$$\min_{\{W, H\} \geq 0} \|X - WH^T\|_F$$

Why?

- Images
- Count Data
- (Hyper) Graphs
- Probabilities



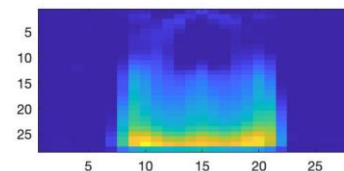
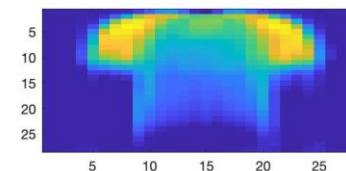
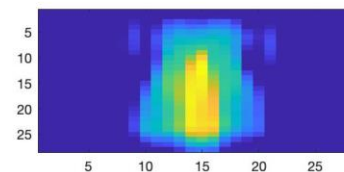
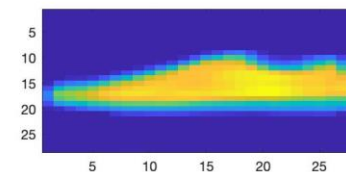
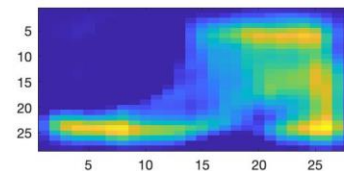
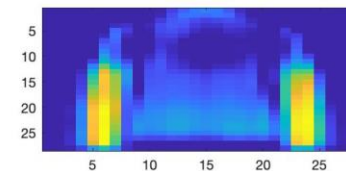
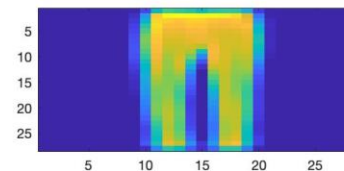
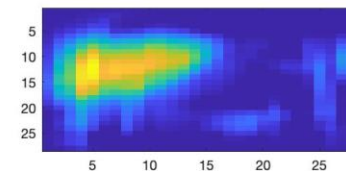
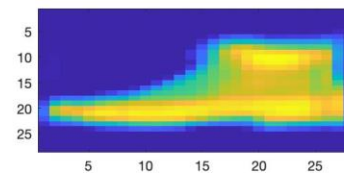
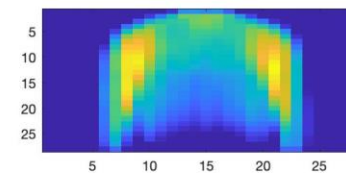
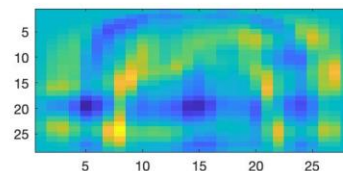
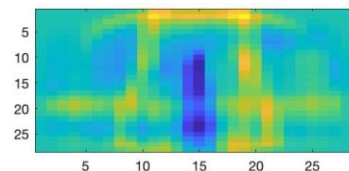
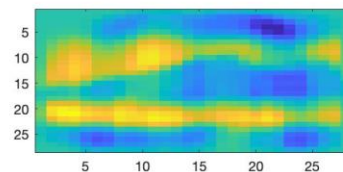
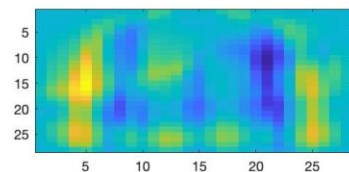
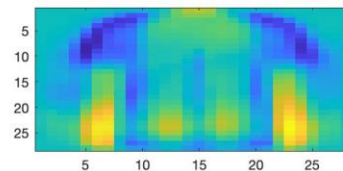
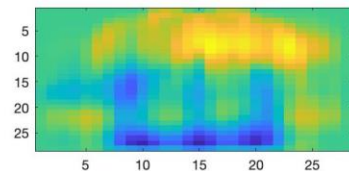
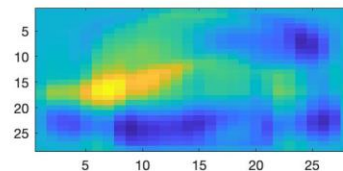
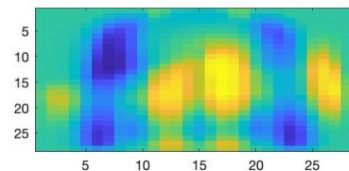
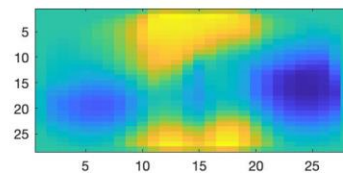
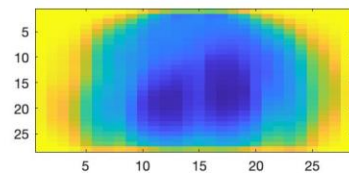
# NMF Example : Fashion MNIST Data Set

Clothing Categories :

- T-shirt/top
- Trouser
- Pullover
- Dress
- Bag

Ect ...

**Interpretability  
via parts-based  
representation**



SVD

NMF

# Netflix Prize

In 2006 Netflix offered \$1 million dollars to improve their recommendation algorithm by **%10**.

In 2009 a team improved the algorithm by **%10.6** and won.

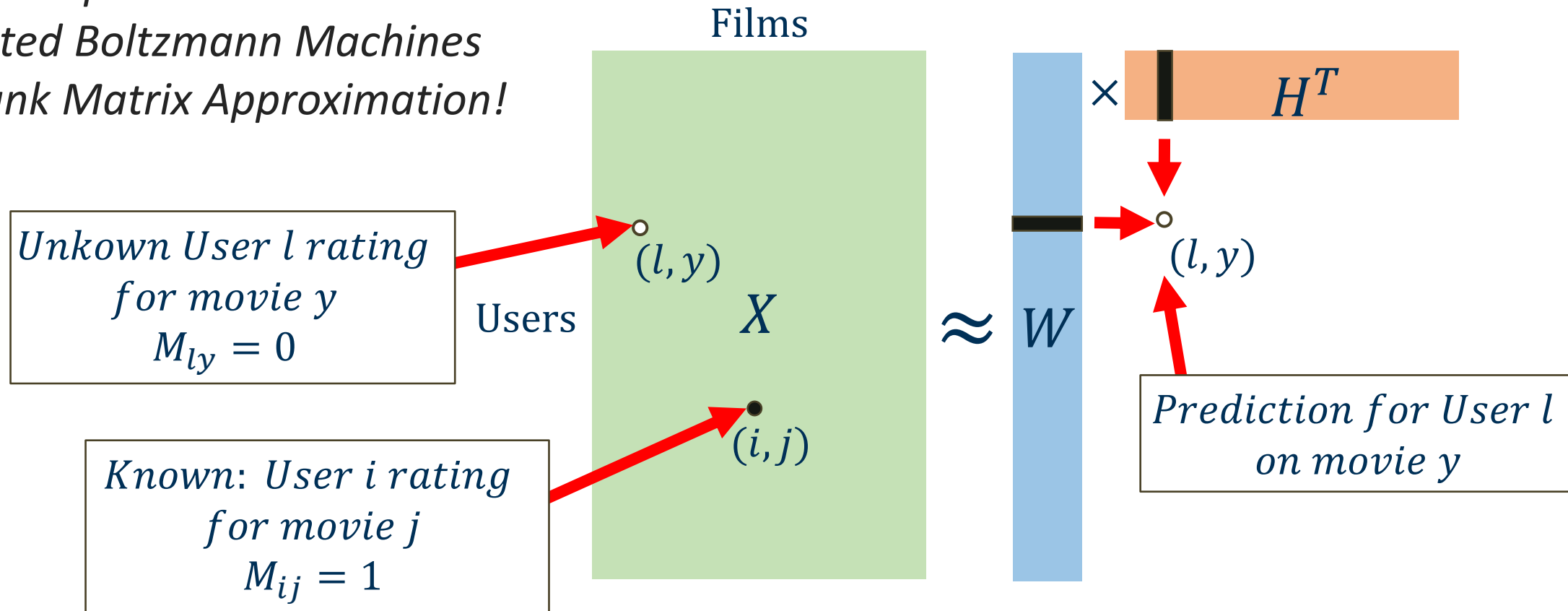
Netflix did not use their algorithm but adopted one that improved their method by only **%8.43...**

**Why?**

# The Adopted Method

2 Main Techniques :

1. Restricted Boltzmann Machines
2. Low-rank Matrix Approximation!



$$\min_{\{W, H\} \geq 0} \|M * (X - WH^T)\|_F$$

# Again... Why?

In 2006 Netflix offered \$1 million dollars to improve their recommendation algorithm by %10.

In 2009 a team improved the algorithm by %10.6 and won.

Netflix did not use their algorithm but adopted one that improved their method by only %8.43...

## Why?

1. Update-ability

2. Scalability

- Training data had 100 million ratings
- Real data had more than 5 billion



# How to Scale

Phillip B. Gibbons @CMU

1. **Scale up** : increase resources on a single node
2. **Scale out** : use multiple nodes
3. **Scale down** : reduce the amount of data or resources needed

<https://github.com/ramkikannan/planc>

S. Eswar, K. Hayashi, G. Ballard, R. Kannan, M. A. Matheson, and H. Park, “Planc: Parallel low-rank approximation with nonnegativity constraints,” ACM Trans. Math. Softw.

S. Eswar, K. Hayashi, G. Ballard, R. Kannan, R. Vuduc, and H. Park, “Distributed-memory parallel symmetric nonnegative matrix factorization,” in SC20: International Conference for High Performance Computing, Networking, Storage and Analysis, 2020

S. Eswar, B. Cobb, K. Hayashi, R. Kannan, G. Ballard, R. Vuduc, and H. Park. 2023. “Distributed-Memory Parallel JointNMF”. In Proceedings of the 37th International Conference on Supercomputing (ICS '23).

# Computing a Nonnegative Matrix Factorization

$$\min_{\{W,H\} \geq 0} \|X - WH^T\|_F$$

Nonlinear – optimization problem!  
Its NP-HARD!

*Repeat Until Converged:*

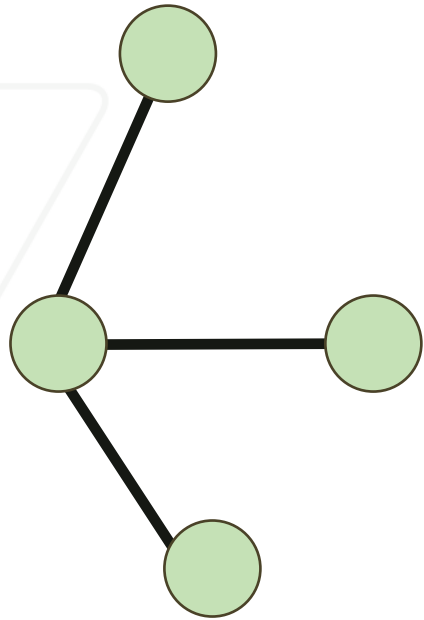
1.  $W^{new} \leftarrow \min_{\{W\} \geq 0} \|X - WH^T\|_F$
2.  $H^{new} \leftarrow \min_{\{H\} \geq 0} \|X - W^{new}H^T\|_F$

Nonnegative  
Least Squares  
Problems

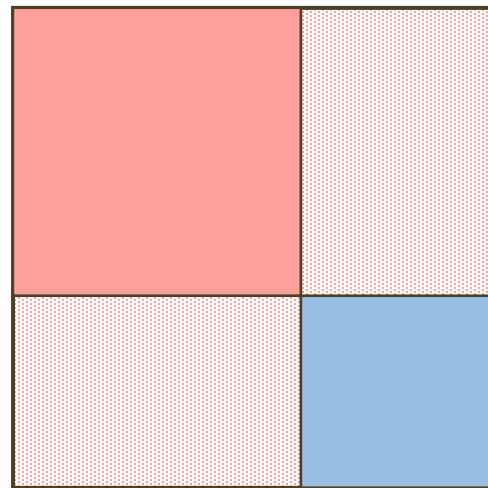
# Symmetric NMF

$$\min_{\{H\} \geq 0} \|A - HH^T\|_F$$
$$A = A^T$$

Paper



Graph Clustering:



$\approx$



+



# Random Compression for SymNMF

1.  $\Omega_{ij} = \mathcal{N}(0,1)$ , a random matrix
2.  $QR = A\Omega$ , compute orthonormal basis
3.  $T = Q^T A Q \leftarrow$  small
4.  $A \approx QTQ^T$

$$A \times \Omega = A\Omega$$

$$A\Omega = Q \times R$$

$Q^T Q = I$

$$\min_{\{H\} \geq 0} \|A - HH^T\|_F \rightarrow \min_{\{H\} \geq 0} \|QTQ^T - HH^T\|_F$$

$QT(Q^T v)$  is faster than  $Av$

$Q$  is an approximate basis for  $A$

# Dense Problem: Hypergraph

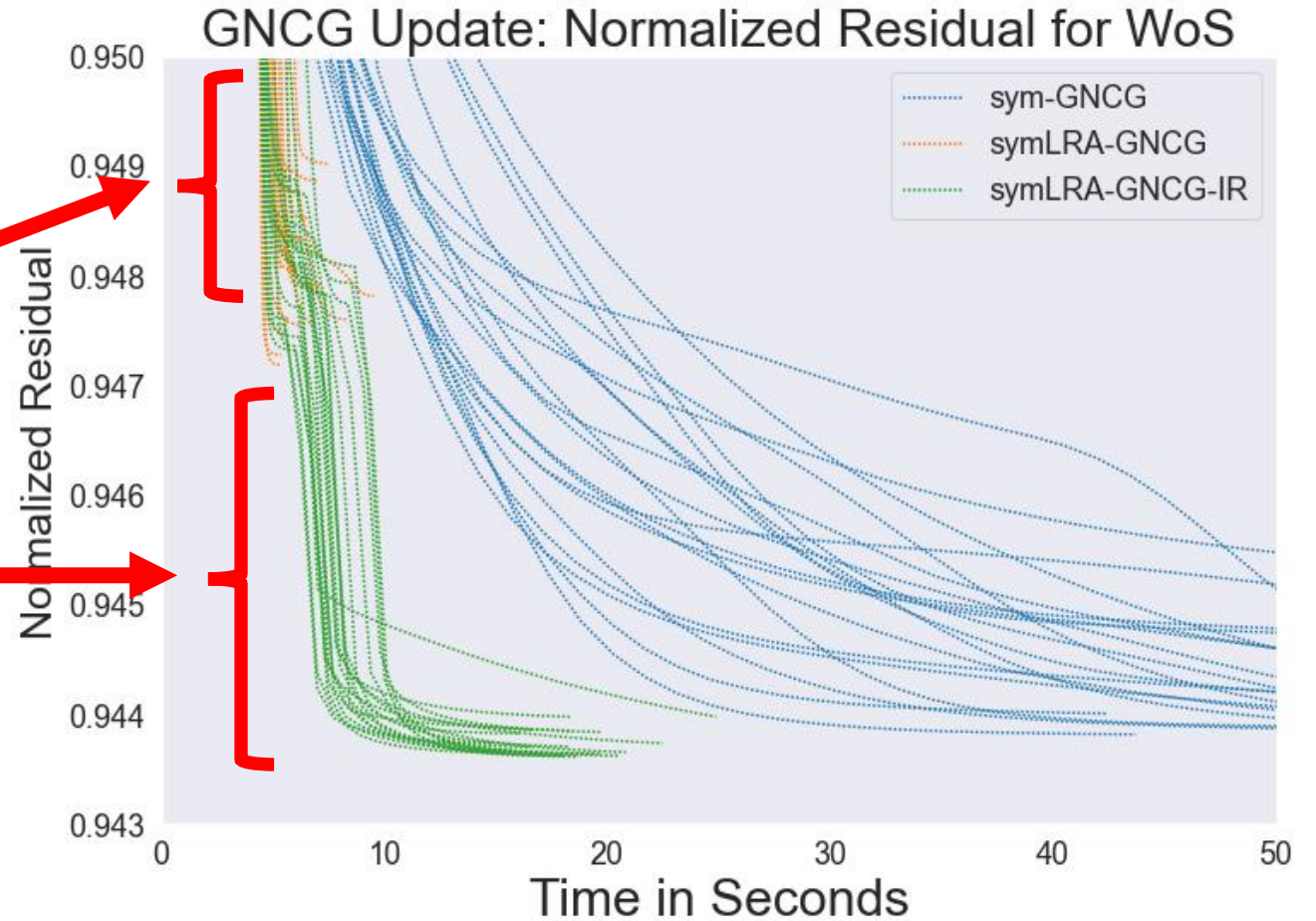
Start with:

$$\min_{\{H\} \geq 0} \|QTQ^T - HH^T\|_F$$

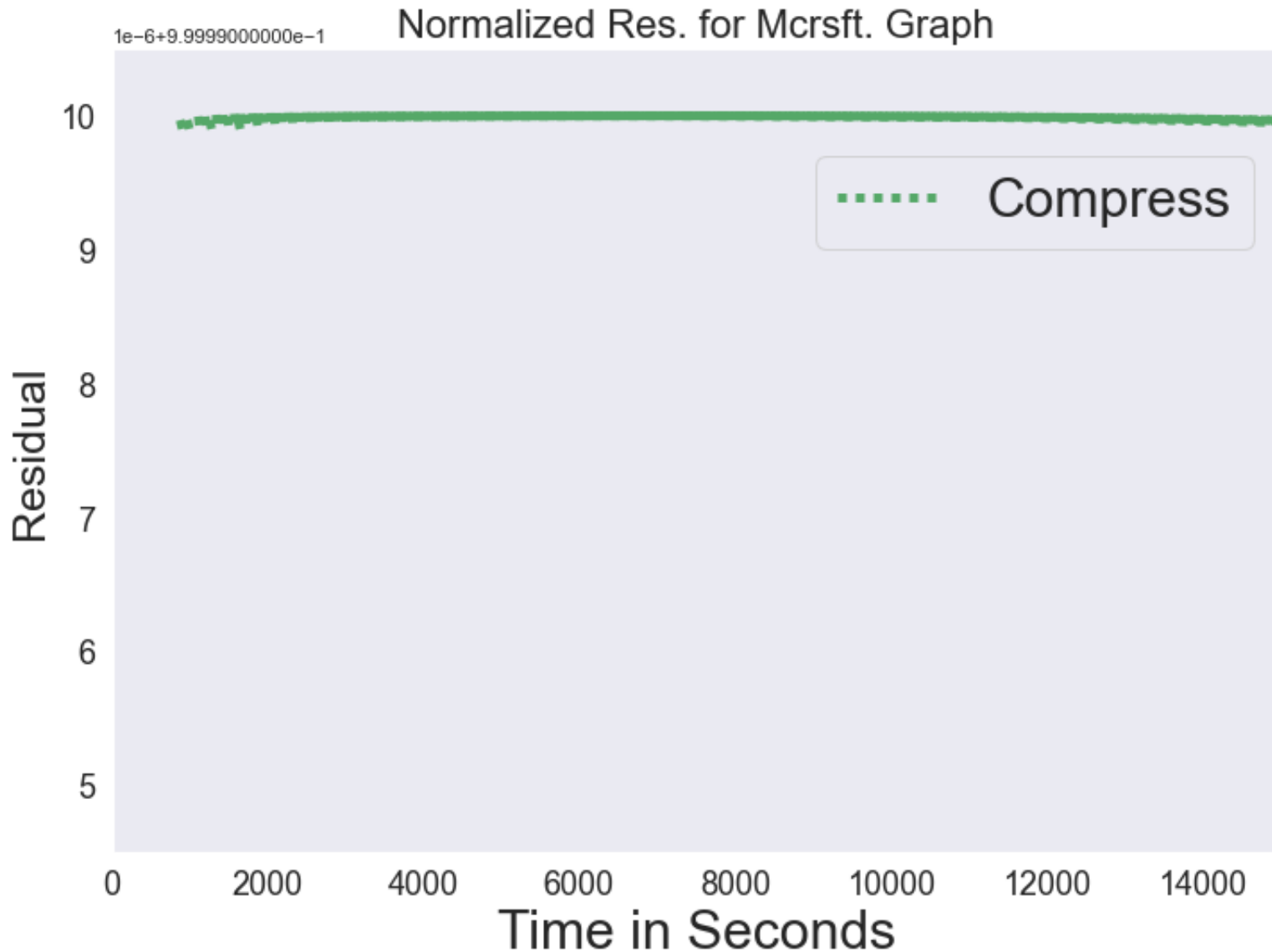
Switch to:

$$\min_{\{H\} \geq 0} \|X - HH^T\|_F$$

- 4x speed up
- Preserve solution quality



# Sparse Problem : Microsoft Academic Graph



Microsoft Academic Graph

1. ~37 million vertices

2. ~ 1 billion edges

# Row Sampling Least Squares

$$\min_x \|Ax - b\|_2^2$$

$A$  is  $m \times k$  and full rank  
 $x$  is length  $k$   
 $b$  is length  $m$

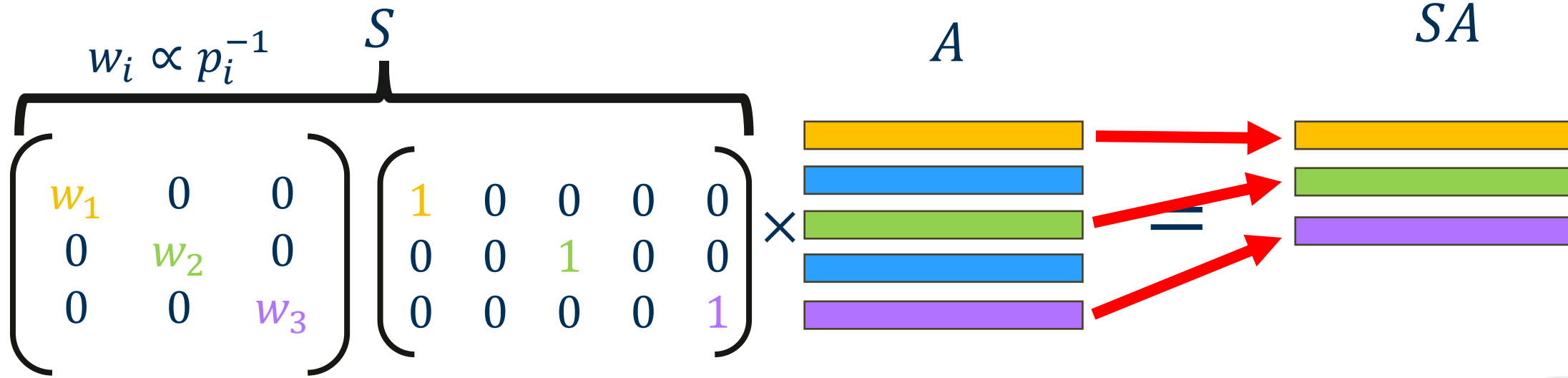


Obtain a distribution  
 $p_1, \dots, p_i, \dots, p_m$   
 For the rows of  $A$  and  $b$

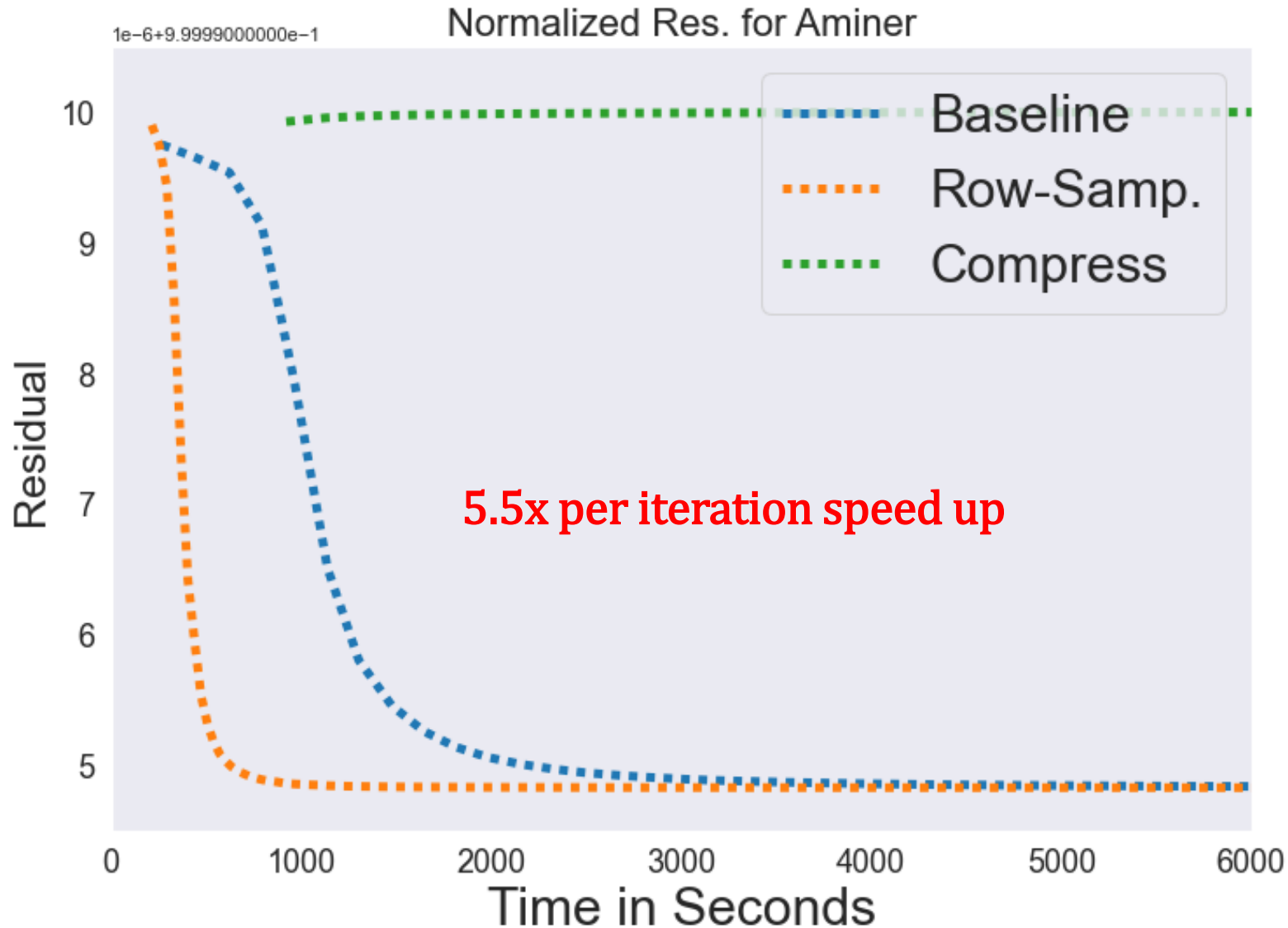


$$\min_x \|SAx - Sb\|_2^2$$

$S$  is  $s \times m, s < m$   
 $S$  samples row of  $A$  and  $b$



# Sparse Problem : Microsoft Academic Graph



Microsoft Academic Graph

1. ~37 million vertices

2. ~ 1 billion edges

**Different types of randomization work for different problem!**



# End



Srinivas Eswar



Ramakrishnan Kannan



Haesun Park



Richard Vuduc



Grey Ballard



Ben Cobb

