

Programming Models for Scientific Computing

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- Not just parallel programming
 - And not just “classical” programming languages – python, Matlab, multi-lingual programs
- (At least) Two goals
 - Clear, maintainable programs
 - “Productivity”
 - Performance
 - Otherwise, you don’t need parallelism
- One more requirement
 - Interoperability with components (library routines) written in other languages
- Most parallel programming models consist of
 - A conventional single-threaded model
 - A parallel coordination layer

Single Threaded Languages

- Fortran, C, C++ (and many others)
 - No intrinsic parallelism
 - Do provide some features for memory hierarchies
- Programming for memory hierarchy
 - These provide some simple tools that can help the compiler produce better-performing code
- In C/C++
 - `const` – Data is not changed
 - `restrict` (pointers) – roughly, data is not accessed with a different pointer
- In Fortran
 - Arguments to routines are *required* to be distinct
 - As if they had C's `restrict` semantics
 - One of the reasons that Fortran is considered easier to optimize than C
 - Fortran provides intent as well (`IN`, `OUT`, `INOUT`). `IN` can help the compiler
- You should *always* use the correct declaration
 - Compilers continue to improve and to exploit this knowledge
 - Compiler may also check whether you told the truth



One more issue - Aligned memory

- Some special features require that operands be aligned on double-word (e.g., 16-byte) boundaries

Example: Using const and restrict

- ```
Void daxpy(int n, const double * restrict x,
 double * restrict y, double a)
{
 register int i;
 for (i=0; i<n; i++)
 y[i] = a * x[i] + y[i];
}
```
- Using const and restrict allows a high-quality compiler to perform optimizations that *are not correct* when const and restrict are not present.
- Don't let experiences with old, bad compilers convince you to ignore these
  - Factor of two benefits are possible



# Illustrating the Programming Models

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- Learning each programming model takes more than an hour 😊
  - This section will show samples of programming models, applied to one simple operation (matrix-vector multiply on a regular grid)
  - For more information, consider
    - Tutorials (e.g., at SC09)
    - Taking a parallel programming class
    - Reading books on different models



# The Poisson Problem

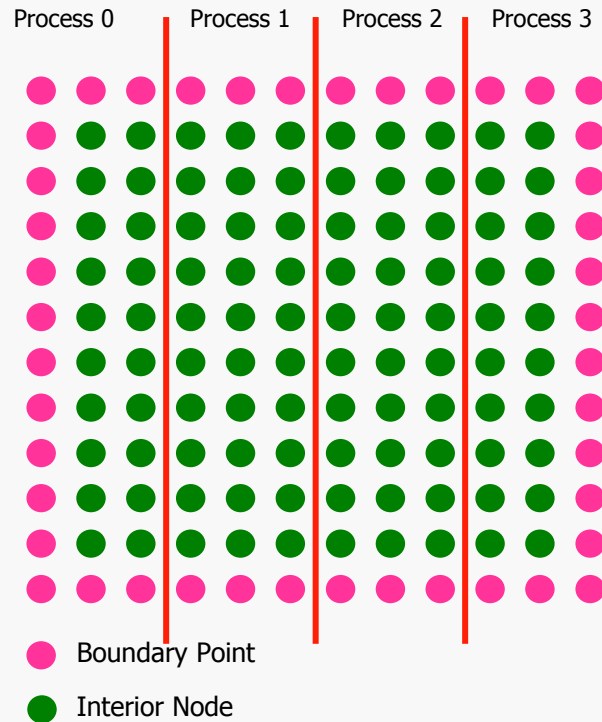
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- Simple elliptic partial differential equation
- Occurs in many physical problems
  - Fluid flow, electrostatics, equilibrium heat flow
- Many algorithms for solution
- We illustrate a sub-optimal one, since it is easy to understand and is typical of a data-parallel algorithm



# Jacobi Iteration (Fortran Ordering)

- Simple parallel data structure



- Processes exchange columns with neighbors
- Local part declared as `xlocal(m,0:n+1)`

```
real u(0:n,0:n), unew(0:n,0:n), f(1:n, 1:n), h
```

```
! Code to initialize f, u(0,*), u(n:*), u(*,0), and
! u(*,n) with g
```

```
h = 1.0 / n
```

```
do k=1, maxiter
```

```
 do j=1, n-1
```

```
 do i=1, n-1
```

```
 unew(i,j) = 0.25 * (u(i+1,j) + u(i-1,j) + &
 u(i,j+1) + u(i,j-1) - &
 h * h * f(i,j))
```

```
 enddo
```

```
 enddo
```

```
! code to check for convergence of unew to u.
```

```
! Make the new value the old value for the next iteration
```

```
u = unew
```

```
enddo
```





# Adding SMP Parallelism

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- Threads
  - A thread is an address space with a program counter.
  - A process contains one or more threads
  - Libraries such as pthreads (for POSIX threads) allow you to create and manipulate new threads
- Threads are a very complex and error-prone programming model.
  - It is better, when possible, to use a higher level SMP programming model...



- OpenMP provides both
  - Loop parallelism
  - Task parallelism
- OpenMP is a set of compiler directives (in comments) and library calls
- The comments direct the execution of loops in parallel in a convenient way.
- Data placement is not controlled, so performance is hard to get except on machines with real shared memory
- There are also subtle memory consistency issues (i.e., when is data written by one thread visible to another)

# OpenMP Version

```
real u(0:n,0:n), unew(0:n,0:n), f(1:n-1, 1:n-1), h
```

```
! Code to initialize f, u(0,*), u(n:*), u(*,0),
! and u(*,n) with g
```

```
h = 1.0 / n
```

```
do k=1, maxiter
```

```
!$omp parallel
```

```
!$omp do
```

```
do j=1, n-1
```

```
do i=1, n-1
```

```
unew(i,j) = 0.25 * (u(i+1,j) + u(i-1,j) + &
u(i,j+1) + u(i,j-1) - &
h * h * f(i,j))
```

```
enddo
```

```
enddo
```

```
!$omp enddo
```

```
! code to check for convergence of unew to u.
```

```
! Make the new value the old value for the next iteration
```

```
u = unew
```

```
!$omp end parallel
```

```
enddo
```





# Distributed Shared Memory Parallelism

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- Three major approaches
- Two-sided message passing (MPI-1)
  - Point-to-point
  - Collective
- One-sided (or Put/Get) models
  - SHMEM, Global Arrays, MPI-2 RMA
- Languages supporting distributed data structures (PGAS)
  - HPF (historical), UPC, CoArray Fortran
  - Future directions: ZPL (historical), Chapel, X10, Fortress



- The Message-Passing Interface (MPI) is a standard library interface specified by the MPI Forum
- It implements the message passing model, in which the sending and receiving of messages combines both data movement and synchronization. Processes have separate address spaces.
- `Send(data, destination, tag, comm)` in one process matches `Receive(data, source, tag, comm)` in another process, at which time data is copied from one address space to another
- Data can be described in many flexible ways
- `SendReceive` can be used for exchange
- Callable from Fortran-77, Fortran-90, C, C++ as specified by the standard
  - Other bindings (Python, java) available, non-standard

# MPI Version – Two Sided

```
use mpi
real u(0:n,js-1:je+1), unew(0:n,js-1:je+1)
real f(1:n-1, js:je), h
integer nbr_down, nbr_up, status(MPI_STATUS_SIZE), ierr
! Code to initialize f, u(0,*), u(n:*), u(*,0), and
! u(*,n) with g
h = 1.0 / n
do k=1, maxiter
 ! Send down
 call MPI_Sendrecv(u(1,js), n-1, MPI_REAL, nbr_down, k &
 u(1,je+1), n-1, MPI_REAL, nbr_up, k, &
 MPI_COMM_WORLD, status, ierr)

 ! Send up
 call MPI_Sendrecv(u(1,je), n-1, MPI_REAL, nbr_up, k+1, &
 u(1,js-1), n-1, MPI_REAL, nbr_down, k+1,&
 MPI_COMM_WORLD, status, ierr)

 do j=js, je
 do i=1, n-1
 unew(i,j) = 0.25 * (u(i+1,j) + u(i-1,j) + u(i,j+1) + u(i,j-1) - &
 h * h * f(i,j))
 enddo
 enddo
 ! code to check for convergence of unew to u.
 ! Make the new value the old value for the next iteration
 u = unew
enddo
```



# MPI Version – One Sided

```
use mpi
real u(0:n,js-1:je+1), unew(0:n,js-1:je+1)
real f(1:n-1, js:je), h
real (kind=MPI_ADDRESS_KIND) off
integer nbr_down, nbr_up, status(MPI_STATUS_SIZE), ierr
! Code to initialize f, u(0,*), u(n:*), u(*,0), and
! u(*,n) with g
! Code to create the MPI memory window
h = 1.0 / n
do k=1, maxiter
 ! Send downcall MPI_Win_fence(win, ierr)
 off = (je-js+2)*(n+1)
 call MPI_Put(u(1,js), n-1, MPI_REAL, nbr_down, off, n-1, MPI_REAL, win, ierr)
 ! Send up
 off = 0
 call MPI_Put(u(1,je), n-1, MPI_REAL, nbr_up, off, n-1, MPI_REAL, win, ierr)
 call MPI_Win_fence(win, ierr)
 do j=js, je
 do i=1, n-1
 unew(i,j) = 0.25 * (u(i+1,j) + u(i-1,j) + u(i,j+1) + u(i,j-1) - &
 h * h * f(i,j))
 enddo
 enddo
 ! code to check for convergence of unew to u.
 ! Make the new value the old value for the next iteration
 u = unew
enddo
```



- HPF is a specification for an extension to Fortran 90 that focuses on describing the distribution of data among processes in structured comments.
- Thus an HPF program is also a valid Fortran-90 program and can be run on a sequential computer
- All communication and synchronization is provided by the compiled code, and hidden from the programmer



```
real u(0:n,0:n), unew(0:n,0:n), f(0:n, 0:n), h
!HPF$ DISTRIBUTE u(:,BLOCK)
!HPF$ ALIGN unew WITH u
!HPF$ ALIGN f WITH u
```

```
! Code to initialize f, u(0,*), u(n:*), u(*,0),
! and u(*,n) with g
```

```
h = 1.0 / n
```

```
do k=1, maxiter
```

```
 unew(1:n-1,1:n-1) = 0.25 * &
 (u(2:n,1:n-1) + u(0:n-2,1:n-1) + &
 u(1:n-1,2:n) + u(1:n-1,0:n-2) - &
 h * h * f(1:n-1,1:n-1))
```

```
! code to check for convergence of unew to u.
```

```
! Make the new value the old value for the next iteration
```

```
 u = unew
enddo
```





# The PGAS Languages

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- PGAS (Partitioned Global Address Space) languages attempt to combine the convenience of the global view of data with awareness of data locality, for performance
  - Co-Array Fortran, an extension to Fortran-90)
  - UPC (Unified Parallel C), an extension to C
  - Titanium, a parallel version of Java



# Co-Array Fortran

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- SPMD – Single program, multiple data
- Replicated to a number of images
- Images have indices 1,2, ...
- Number of images fixed during execution
- Each image has its own set of local variables
- Images execute asynchronously except when explicitly synchronized
- Variables declared as co-arrays are accessible of another image through set of array subscripts, delimited by [ ] and mapped to image indices by the usual rule
- Intrinsic: `this_image`, `num_images`, `sync_all`, `sync_team`, `flush_memory`, collectives such as `co_sum`
- A version of Co\_Array Fortran is a part of the Fortran 2008 draft standard (mostly a subset of the full CAF)



```
real u(0:n,js-1:je+1,0:1)[*], f (0:n,js:je), h
integer np, myid, old, new
np = NUM_IMAGES()
myid = THIS_IMAGE()
new = 1
old = 1-new
! Code to initialize f, and the first and last columns of u on the extreme
! processors and the first and last row of u on all processors
h = 1.0 / n
do k=1, maxiter
 if (myid .lt. np) u(:,js-1,old)[myid+1] = u(:,je,old)
 if (myid .gt. 0) u(:,je+1,old)[myid-1] = u(:,js,old)
 call sync_all
 do j=js,je
 do i=1, n-1
 u(i,j,new) = 0.25 * (u(i+1,j,old) + u(i-1,j,old) + u(i,j+1,old) + u(i,j-1,old) - &
 h * h * f(i,j))
 enddo
 enddo
 ! code to check for convergence of u(:,:,new) to u(:,:,old).
 ! Make the new value the old value for the next iteration
 new = old
 old = 1-new
enddo
```

- UPC is an extension of C (not C++) with shared and local pointers
- Provides distributed arrays
  - “Global view” simplifies programming
- Provides a “Put-Get” model of programming
  - You can dereference a shared pointer
- No standard, but an active community and multiple implementations, including some “vendor” implementations

```
#include <upc.h>
#define n 1024
shared [*] double u[n+1][n+1];
shared [*] double unew[n+1][n+1];
shared [*] double f[n][n];
int main() {
 int maxiter = 100;
 // Code to initialize f, u(0,*), u(n:*), u(*,0), and
 // u(*,n) with g
 double h = 1.0 / n;
 for (int k=0; k < maxiter; k++) {
 for (int i=1; i < n; i++) {
 upc_forall (int j=1; j < n; j++; &unew[i][j]) {
 unew[i][j] = 0.25 * (u[i+1][j] + u[i-1][j] + u[i][j+1] + u[i][j-1] -
 h * h * f[i][j]);
 }
 }
 upc_barrier;
 // code to check for convergence of unew to u.
 // Make the new value the old value for the next iteration
 for (int i = 1; i < n; i++) {
 upc_forall(int j = 1; j < n; j++; &u[i][j]) {
 u[i][j] = unew[i][j];
 }
 }
 }
}
```



# Global Operations

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- Some operations involve all processes
  - Or a collection, such as all processes in a “row” or “column”
- Example: checking for convergence
  - Need to compute the two-norm of the difference between successive iterations for the *entire data* array



# Serial Version

---

```
real u(0:n,0:n), unew(0:n,0:n), twonorm
```

```
! ...
```

```
twonorm = 0.0
```

```
do j=1, n-1
```

```
 do i=1, n-1
```

```
 twonorm = twonorm + (unew(i,j) - u(i,j))**2
```

```
 enddo
```

```
enddo
```

```
twonorm = sqrt(twonorm)
```

```
if (twonorm .le. tol) ! ... declare convergence
```



```
real u(0:n,0:n), unew(0:n,0:n), twonorm
```

```
! ..
```

```
 twonorm = 0.0
```

```
!$omp parallel
```

```
!$omp do private(ldiff) reduction(+:twonorm)
```

```
 do j=1, n-1
```

```
 do i=1, n-1
```

```
 ldiff = (unew(i,j) - u(i,j))**2
```

```
 twonorm = twonorm + ldiff
```

```
 enddo
```

```
 enddo
```

```
!$omp enddo
```

```
!$omp end parallel
```

```
 twonorm = sqrt(twonorm)
```

```
enddo
```



```
use mpi
real u(0:n,js-1:je+1), unew(0:n,js-1:je+1), twonorm
integer ierr
```

```
! ...
```

```
twonorm_local = 0.0
do j=js, je
 do i=1, n-1
 twonorm_local = twonorm_local + &
 (unew(i,j) - u(i,j))**2
 enddo
enddo
call MPI_Allreduce(twonorm_local, twonorm, 1, &
 MPI_REAL, MPI_SUMM, MPI_COMM_WORLD, ierr)
twonorm = sqrt(twonorm)
if (twonorm .le. tol) ! ... declare convergence
```



```
real u(0:n,0:n), unew(0:n,0:n), twonorm
!HPF$ DISTRIBUTE u(:,BLOCK)
!HPF$ ALIGN unew with u
!HPF$ ALIGN f with u

! ...

twonorm = sqrt (&
 sum ((unew(1:n-1,1:n-1) - u(1:n-1,1:n-1))**2))
if (twonorm .le. tol) ! ... declare convergence
enddo
```





# The HPCS languages

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- DARPA funded three vendors to develop next-generation languages for programming next-generation petaflops computers
  - Fortress (Sun)  
<http://projectfortress.sun.com/Projects/Community>
  - X10 (IBM) <http://x10-lang.org/>
  - Chapel (Cray) <http://chapel.cs.washington.edu/>
- All are global-view languages, but also with some notion for expressing locality, for performance reasons.
  - They are more abstract than UPC and CAF in that they do not have a fixed number of processes.





# Other Programming Approaches

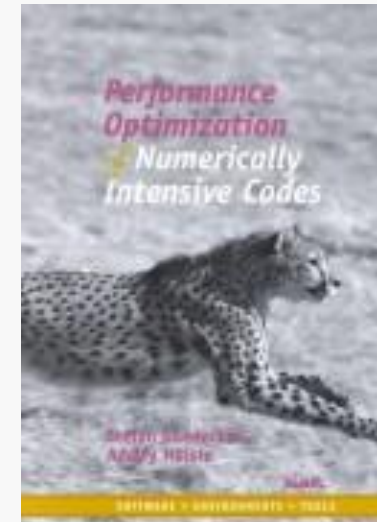
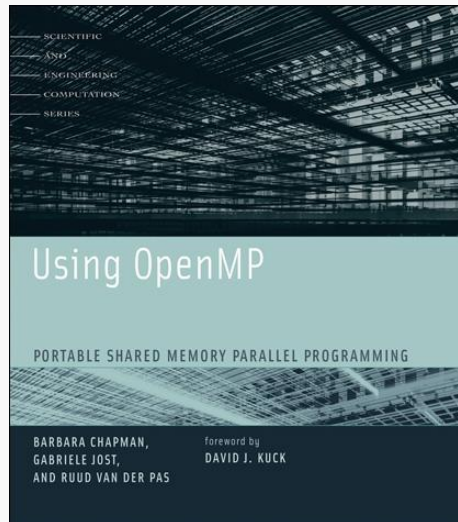
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- Programming accelerators (GPGPUs)
  - Emerging (but low-level) languages such as CUDA, OpenCL
- Integrated Development Environments (Eclipse)
  - Powerful but with a high learning curve
- Automating code optimization and tuning (Annotations, Autotuning)
  - A number of research tools, particularly at the DOE Labs, but not generally ready for average users



- Short Term Projects
  - Use whatever best fits your application
    - This might be libraries that provide all the needed methods
  - CAF good for regular data structures; UPC also has some good support
  - If you use a PGAS language, make sure
    - You have access to CAF or UPC
    - The environments produce efficient code
- Long Term Projects
  - MPI, possibly mixed with OpenMP, remains the reliable and efficient standard
  - But make use of existing software components as much as possible

- Using MPI
- Using OpenMP
- Performance Optimization of Numerically Intensive Codes



- Scalable and Efficient Algorithms
  - Combines applied mathematics with architecture influences
    - Picking algorithms with maximum FLOPS often wrong
    - Picking algorithms with minimum FLOPs often wrong
  - David Keyes presents an overview of HPC algorithms next